

# Reducing tandem queues (sketch)

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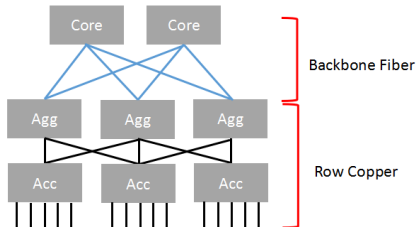
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# Data centres and tandem queues

[Introduction to data centres and the important questions surrounding their functionality]

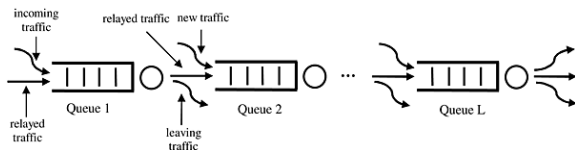
## Data centres and tandem queues

A data centre might be modelled as a (tiered) network of queues. The customers are packets of data, which arrive from a host, travel through the aggregating layers, and are then sent to another host.



# Data centres and tandem queues

Let's focus on an **single end-to-end route** through the network - this can be modelled as a sequence of queues in tandem, with incoming and outgoing cross-traffic.



The inner workings of this tandem system are complicated - we might wish to reproduce its behaviour (ie the sojourn or holding times of the customers) more simply).

## Data centres and tandem queues

One way to do this is to see if we can replicate the behaviour using a shorter tandem system, the simplest case being a single queue:



This is still quite a general set-up, with totally arbitrary cross-traffic and service distributions, so task does not seem unreasonable.

# Analytic results

- ▶ One way to approach the problem is to assume a fully specified probabilistic model - arrival, cross-traffic, and service distributions, scheduling disciplines, buffer sizes and so forth
- ▶ There is some nice existing work on reducing networks of queues to smaller systems

# Reduction Methods for Tandem Queuing Systems (Friedman 1963)

- ▶ Tandem queue  $(Q_1, \dots, Q_n)$
- ▶ General arrival process,
- ▶  $m_i$  servers in queue  $i$ , each serving at constant rate  $s_i$
- ▶ No cross-traffic, infinite buffers
- ▶ First note that the queue order in this system doesn't affect the exit distribution of the customers
- ▶ Say  $Q_i$  **dominates**  $Q_j$  if in the (different!) system  
→  $Q_i \rightarrow Q_j$ , no customer waits at  $Q_j$  for any input process to  $Q_i$

# Reduction Methods for Tandem Queuing Systems (Friedman)

- ▶ There is a very simple criterion for domination in our constant service queues, namely that  $s_j \leq \lceil m_j/m_i \rceil s_i$
- ▶ The property of dominance is **transient** and **persistent** (ie it operates across interposed queues)
- ▶ Then, for example, in a tandem system in which one queue dominates all the others, rearrange to bring that queue to the front
- ▶ Now no waiting occurs downstream, so all other queues can be replaced by a constant delay



# Reduction Methods for Tandem Queuing Systems (Friedman)

- ▶ In general, reduction proceeds as follows:
- ▶ Reorder the  $Q_i$  in ascending order of  $m_i$ , and sub-order by  $s_i$  descending
- ▶ For each  $m_i$ , the queue with the highest  $s_i$  dominates all the others - move these to back of system
- ▶ For remaining queues, apply dominance test and move all dominated queues to back of system
- ▶ Finally, replace the dominated subsystem by a constant delay

# Reduction Methods for Tandem Queuing Systems (Friedman)

Using the notation  $(s_i, m_i)$  for  $Q_i$ , consider as an example:

$$\rightarrow (5; 2) \rightarrow (4; 2) \rightarrow (2; 1) \rightarrow (3; 1) \rightarrow (7; 3) \rightarrow (3; 1)$$

$$\rightarrow (3; 1) \rightarrow (3; 1) \rightarrow (2; 1) \rightarrow (5; 2) \rightarrow (4; 2) \rightarrow (7; 3)$$

$$\rightarrow (3; 1) \rightarrow (5; 2) \rightarrow (7; 3) \rightarrow (3; 1) \rightarrow (2; 1) \rightarrow (4; 2)$$

$$\rightarrow (3; 1) \rightarrow (5; 2) \rightarrow (7; 3) \rightarrow [9]$$

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$$\rightarrow (3; 1) \rightarrow [21]$$

- ▶ It is well-known that  $M/1$  queues in tandem are also interchangeable
- ▶ It may be tempting to weaken the notion of dominance to a probabilistic one, and apply similar ideas - but persistence is not all clear in this case
- ▶ In a data centre, service times are very fast, and queueing behaviour is generated by 'colliding cross-traffic'

## Further analytic results

Further discussion of other queue reduction results, eg:

- ▶ A Tandem Network of Queues with Deterministic Service and Intermediate Arrivals (Shalman and Kaplan 1984)
- ▶ Delay analysis for tandem networks (Konheim and Reiser 1977)
- ▶ Reduction of a polling network to a single node (Beekhuizen, Denteneer, Resing 2008)

# Limitations

- ▶ Fairly strong assumptions are needed on the nature of the arrival processes, service processes and other attributes of the system to make analysis tractable
- ▶ Tend to be exact, but we could be satisfied with approximations
- ▶ Generally we don't have complete information about the inner workings of a data centre
- ▶ Often limited to preserving only first moments - more accuracy is needed to produce good simulations

## Turning the problem around

- ▶ Rather than assuming a fully specified probabilistic model, suppose instead we gather data describing the packet sojourn times through the tandem system
- ▶ Can we use this to construct a single queue with the desired properties?
- ▶ In other words, assuming the results came from a single queue, what can we infer about that queue?

# Inverse problems in queueing

- ▶ Inference by probing
- ▶ The M/G/1 queue
- ▶ Lindsley's equation

# The G/G/1 queue without cross-traffic

- ▶ The G/G/1 queue is the most natural simple model with which to begin
- ▶ Suppose we have access to a set of arrival and departure times for the customers
- ▶ Then the service times are fully specified
- ▶ How good/poor is the approximation that these are independent samples from some distribution?



# The G/G/1 queue with cross-traffic

- ▶ Discussion of several methods of determining cross-traffic and service distributions, with numerical examples, will make up the rest of the talk